Implementation of a Mortar Mixed Finite Element Method using a Multiscale Flux Basis

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Abstract
This paper provides a new implementation of a multiscale mortar mixed finite element method for second order elliptic problems. The algorithm uses non-overlapping domain decomposition to reformulate a fine scale problem as a coarse scale mortar interface problem, which is then solved using an iterative method. The original implementation by Arbogast, Pencheva, Wheeler, and Yotov, Multiscale Model. Simul. 2007, required solving one local Dirichlet problem on each subdomain per interface iteration. We alter this implementation by forming a multiscale flux basis. This basis consists of mortar functions representing the individual flux responses for each mortar degree of freedom, on each subdomain independently. The computation of these basis functions requires solving a fixed number of Dirichlet subdomain problems. Taking linear combinations of the multiscale flux basis functions replaces the need to solve any Dirichlet subdomain problems during the interface iteration. This new implementation yields the same solution as the original implementation, and is computationally more efficient in cases where the number of interface iterations is greater than the number of mortar degrees of freedom per subdomain. The gain in computational efficiency increases with the number of subdomains.

Keywords. Multiscale, Mortar Finite Element, Mixed Finite Element, Porous Media Flow

1 Introduction
This paper provides a new way of implementing the multiscale mortar mixed finite element method (MMMFEM) which was proposed by Arbogast, Pencheva, Wheeler,

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and Yotov in 2007 [7]. We consider a second order elliptic equation (1), which models single phase flow in porous media. The permeability tensor $K$ varies on a fine scale and so do the velocity $u$ and the pressure $p$. Resolving the solution on the fine scale is often computationally infeasible, necessitating multiscale approximations. Our choice of mixed finite element method for the discretization is motivated by its local element-wise conservation of mass and accurate velocity approximation.

The MMMFEM was proposed in [7] as an alternative to existing multiscale methods, such as the variational multiscale method [18, 19, 4, 2, 5, 3] and multiscale finite elements [16, 17, 14, 11, 20, 1]. The latter two approaches are closely related [5]. In all three methods the domain is decomposed into a series of small subdomains (coarse grid) and the solution is resolved globally on the coarse grid and locally (on each coarse element) on a fine grid. All three are based on a divide and conquer approach: solving relatively small fine scale subdomain problems that are only coupled together through a reduced number (coarse scale) degrees of freedom.

The variational multiscale method and multiscale finite elements both compute a multiscale basis by solving local fine scale problems with boundary conditions or a source term corresponding to the coarse scale degrees of freedom. This basis is then used to solve the coarse scale problem. The MMMFEM uses a non-overlapping domain decomposition algorithm which introduces a Lagrange multiplier space on the subdomain interfaces to weakly impose certain continuity conditions. By eliminating the subdomain unknowns the global fine scale problem is reduced to an interface problem, which is solved using an iterative method. The domain decomposition algorithm was originally developed for the case of matching grids [15] and then extended to the case of non-matching grids using mortar finite elements [28, 6]. This generalization allows for extremely flexible finite element partitions, as both the fine scale elements across subdomain interfaces and the subdomains themselves (i.e., the coarse grid) may be spatially non-conforming. Moreover, one has the ability to vary the interface degrees of freedom [26, 23, 7]. If only a single mortar grid element is used per interface, the resulting approximation is comparable to the one in the variational multiscale method or multiscale finite elements. In the MMMFEM framework, a posteriori error estimators [27] can be employed to adaptively refine the mortar grids where necessary to improve the global accuracy. Furthermore, higher order mortar approximation can be used to compensate for the coarseness of the mortar grid and obtain fine scale convergence of the error [7]. Thus, the MMMFEM is more flexible than the variational multiscale method and multiscale finite elements. Another observation is that the MMMFEM resolves the flux through the coarse interfaces on the fine scale, which is not the case for the other two approaches.

The original implementation of the MMMFEM in [7] requires solving one Dirichlet fine scale subdomain problem per interface iteration. As a result the number of subdomain solves increases whit the dimension of the coarse space, making it difficult to compare the computational efficiency of the method to other existing multiscale methods. In this paper we alter this implementation by forming what we call a
Multiscale Flux Basis, before the interface iteration begins. This basis consists of mortar functions representing the individual flux responses from each mortar degree of freedom, on each subdomain independently. These basis functions may also be described as traces of the discrete Green’s functions corresponding to the mortar degrees of freedom along the subdomain interfaces. The computation of these basis functions requires solving a fixed number of Dirichlet subdomain problems. Taking linear combinations of the Multiscale Flux Basis functions replaces the need to solve any Dirichlet subdomain problems during the interface iteration. This new implementation yields the same solution as the original implementation and makes the MMMFEM comparable to the variational multiscale method and multiscale finite elements in terms of computational efficiency. In our numerical experiments we compare the computational cost of the new implementation to the one for the original implementation with and without preconditioning of the interface problem. If no preconditioning is used, the Multiscale Flux Basis implementation is computationally more efficient in cases where the number of mortar degrees of freedom per subdomain is less than the number of interface iterations. If balancing preconditioning is used [13, 22], the number of iterations is reduced, but each interface iteration requires three subdomain solves. In this case the Multiscale Flux Basis implementation is more efficient if the number of mortar degrees of freedom per subdomain is less than three times the number of interface iterations.

The format of the paper is as follows. Section 2 introduces the MMMFEM method and its step-by-step formulation leading to its original implementation. Section 3 describes our new implementation. In particular, it introduces the concept of a Multiscale Flux Basis, explains how it is used in the interface iteration, and discusses specific implementation details. Section 4 provides several numerical examples which illustrate the computational efficiency of the Multiscale Flux Basis implementation. Section 5 contains concluding remarks and directions for further work.

2 The MMMFEM

Our model problem is a second-order linear elliptic equation written as a first order system in mixed form, arising in applications to single phase incompressible flow in porous media. The pressure \( p \) and the Darcy velocity \( u \) satisfy the system

\[
\begin{align*}
\alpha p + \nabla \cdot u &= f \quad \text{in } \Omega, \\
\quad u &= -K \nabla p \quad \text{in } \Omega, \\
\quad p &= g_D \quad \text{on } \Gamma_D, \\
\quad u \cdot n &= g_N \quad \text{on } \Gamma_N.
\end{align*}
\]

Here \( \Omega \) is a bounded domain in \( \mathbb{R}^d \) (\( d = 2 \) or 3) with boundary \( \partial \Omega = \Gamma_D \cup \Gamma_N, \Gamma_D \cap \Gamma_N = \emptyset \), and outer unit normal \( n \), \( \alpha(x) \geq 0 \), and \( K(x) \) is a symmetric and uniformly
positive definite permeability tensor with components in $L^\infty(\Omega)$. We assume that $f \in L^2(\Omega)$, $g_D \in H^{1/2}(\Gamma_D)$, and $g_N \in L^2(\Gamma_N)$.

Throughout the paper, $C$ denotes a generic positive constant independent of the discretization parameters $h$ and $H$. For a domain $G \subset \mathbb{R}^d$, the $L^2(G)$ inner product and norm for scalar and vector valued functions are denoted $(\cdot, \cdot)_G$ and $\| \cdot \|_G$, respectively. We omit $G$ in the subscript if $G = \Omega$. For a section of the domain or element boundary $S \subset \mathbb{R}^{d-1}$ we write $(\cdot, \cdot)_S$ and $\| \cdot \|_S$ for the $L^2(S)$ inner product (or duality pairing) and norm, respectively.

### 2.1 Domain Decomposition

The first step in formulating the MMMFEM is to use the domain decomposition approach described in [15] to restrict the model problem into smaller pieces. The domain $\Omega$ is divided into non-overlapping subdomains $\Omega_i$, $i = 1, \ldots, n$, that are allowed to be spatially non-conforming, so we have $\Omega = \bigcup_{i=1}^n \Omega_i$ and $\Omega_i \cap \Omega_j = \emptyset$ for $i \neq j$. Denote the single interface between subdomains $\Omega_i$ and $\Omega_j$ by $\Gamma_{i,j} = \partial \Omega_i \cap \partial \Omega_j$, all interfaces that touch subdomain $\Omega_i$ by $\Gamma_i = \partial \Omega_i \setminus \partial \Omega$, and the union of all interfaces by $\Gamma = \bigcup_{i\neq j} \Gamma_{i,j}$. The domain decomposition can be viewed as a coarse grid on $\Omega$.

System (1) holds within each subdomain $\Omega_i$. The pressure and the normal components of the velocity must be continuous across the interfaces. Equivalently, we seek $(u_i, p_i)$ such that for $i = 1, \ldots, n$,

\begin{align}
\alpha p_i + \nabla \cdot u_i &= f \quad \text{in } \Omega_i, \quad (2a) \\
u_i &= -K \nabla p_i \quad \text{in } \Omega_i, \quad (2b) \\
p_i &= g_D \quad \text{on } \partial \Omega_i \cap \Gamma_D, \quad (2c) \\
u_i \cdot n_i &= g_N \quad \text{on } \partial \Omega_i \cap \Gamma_N, \quad (2d) \\
p_i &= p_j \quad \text{on } \Gamma_{i,j}, \quad i \neq j, \quad (2e) \\
u_i \cdot n_i + u_j \cdot n_j &= 0 \quad \text{on } \Gamma_{i,j}, \quad i \neq j, \quad (2f)
\end{align}

where $n_i$ is the outer unit normal to $\partial \Omega_i$.

### 2.2 Variational Formulation

The weak pressure and velocity spaces for the global problem (1) are

\[ \tilde{W} = L^2(\Omega), \quad \tilde{V} = H(\text{div}; \Omega), \quad \tilde{V}^\gamma = \{ v \in \tilde{V} | v \cdot n = \gamma \text{ on } \Gamma_N \}, \]

where

\[ H(\text{div}; \Omega) = \{ v \in (L^2(\Omega))^d | \nabla \cdot v \in L^2(\Omega) \}. \]

The corresponding dual mixed variational formulation is to find $u \in \tilde{V}^{g_N}$ and $p \in \tilde{W}$ such that

\begin{align}
(K^{-1}u, v)_\Omega - (p, \nabla \cdot v)_\Omega &= -\langle v \cdot n, g_D \rangle_{\partial \Gamma_D} \quad \forall v \in \tilde{V}^0, \quad (3a) \\
(\alpha p, w)_\Omega + (\nabla \cdot u, w)_\Omega &= (f, w)_\Omega \quad \forall w \in \tilde{W}. \quad (3b)
\end{align}
Note that the Neumann boundary condition is imposed essentially in the space \( \tilde{V}^{gN} \), which is different from the weak velocity test space \( \tilde{V}^0 \).

Similarly, define the weak spaces for each subdomain \( \Omega_i \) by

\[
W_i = L^2(\Omega_i), \quad V_i = H(\text{div}; \Omega_i), \quad V^\gamma_i = \{ v \in V_i | v \cdot n = \gamma \text{ on } \partial \Omega_i \cap \Gamma_N \}.
\]

The global weak spaces for the domain decomposition problem (2) are

\[
W = \bigoplus_{i=1}^n W_i, \quad V^\gamma = \bigoplus_{i=1}^n V^\gamma_i.
\]

Note that no continuity is imposed across the interfaces for functions in \( V \) and \( W \).

On the interfaces \( \Gamma \) we introduce a Lagrange multiplier space that has a physical meaning of pressure and is used to weakly impose continuity of the normal velocities:

\[
M = \{ \mu \in H^{1/2}(\Gamma) | \mu|_{\Gamma_i} \in (V_i \cdot n_i)^*, i = 1, \ldots, n \},
\]

where \((\cdot)^*\) denotes the dual space.

The corresponding mixed variational formulation is to find \( u \in V^{gN}, p \in W \), and \( \lambda \in M \) such that for \( i = 1, \ldots, n \),

\[
(K^{-1}u, v)_{\Omega_i} - (p, \nabla \cdot v)_{\Omega_i} = -\langle v \cdot n_i, \lambda \rangle_{\partial \Omega_i \cap \Gamma} - \langle v \cdot n_i, g_D \rangle_{\partial \Omega_i \cap \Gamma_D} \quad \forall v \in V^0_i, \quad (4a)
\]

\[
(\alpha p, w)_{\Omega_i} + (\nabla \cdot u, w)_{\Omega_i} = (f, w)_{\Omega_i} \quad \forall w \in W_i, \quad (4b)
\]

\[
\sum_{i=1}^n \langle u_i \cdot n_i, \mu \rangle_{\Gamma_i} = 0 \quad \forall \mu \in M, \quad (4c)
\]

Since \( V^\gamma \neq \tilde{V}^\gamma \), the extra condition (4c) is needed to weakly enforce the flux continuity lost across the interfaces in the domain decomposition.

The following equivalence result is easy to show.

**Lemma 2.1.** If the solution \((u, p)\) to (3) satisfies (1) in a distributional sense, then \((u, p, p|\Gamma)\) solves (4). Conversely, if \((u, p, \lambda)\) solves (4), then \((u, p)\) solves (3).

### 2.3 Discrete Formulation

The multiscale approach to the mortar mixed finite element method combines a local fine scale discretization within each subdomain with a global coarse scale discretization across subdomain interfaces.

First, independently partition each subdomain \( \Omega_i \) into its own local \( d \)-dimensional quasi-uniform affine mesh \( T_{h_i} \). The faces (or edges) of these meshes are spatially conforming within each subdomain, but are allowed to be non-conforming along subdomain interfaces. Let the maximal element diameter of this fine mesh be \( h_i \), and let the global characteristic fine scale diameter be \( h = \max_{i=1}^n h_i \). Denote the global
fine mesh by $T_h = \bigcup_{i=1}^n T_{h,i}$. Let $V_{h,i} \times W_{h,i} \subset V_i \times W_i$ be a mixed finite element space on the mesh $T_{h,i}$ such that $V_{h,i}$ contains piecewise polynomials of degree $k$ and $W_{h,i}$ contains piecewise polynomials of degree $l$. Examples of such spaces are the RT spaces [25, 21], the BDM spaces [10], the BDFM spaces [9], the BDDF spaces [8], or the CD spaces [12]. Globally, the discrete pressure and velocity spaces for this method are $W_h = \bigoplus_{i=1}^n W_{h,i}$ and $V_h = \bigoplus_{i=1}^n V_{h,i}$.

Second, we partition each subdomain interface $\Gamma_{i,j}$ with a $(d - 1)$-dimensional quasi-uniform affine mesh denoted $T_{H,i,j}$. This mesh will be the mortar space that weakly enforces continuity of normal fluxes for the discrete velocities across the non-matching grids. Let the maximal element diameter of this coarse mesh be $H_{i,j}$, and let the global characteristic coarse scale diameter be $H = \max_{1 \leq i < j \leq n} H_{i,j}$. Denote the global coarse mesh by $T_H = \bigcup_{1 \leq i < j \leq n} T_{H,i,j}$. Let $M_{H,i,j} \subset L^2(\Gamma_{i,j})$ be the mortar space containing continuous or discontinuous piecewise polynomials of degree $m$ where $m \geq k + 1$. Globally, the mortar space for this method is $M_H = \bigoplus_{1 \leq i < j \leq n} M_{H,i,j}$. Notice that this is a nonconforming approximation, as $M_H \not\subset W$. With these finite dimensional subspaces, the multiscale mortar mixed finite element approximation of (4) is to find $u_h \in V^{\mathbb{N}}_h$, $p_h \in W_h$, and $\lambda_H \in M_H$ such that for $i = 1, \ldots, n$,

$$(K^{-1}u_h, v)_{\Omega_i} - (p_h, \nabla \cdot v)_{\Omega_i} = -\langle v \cdot n_i, \lambda_H \rangle_{\partial \Omega_i \cap \Gamma} - \langle v \cdot n_i, g_D \rangle_{\partial \Omega_i \cap \Gamma_D} \forall v \in V^{\mathbb{N}}_h, (5a)$$

$$(\alpha p_h, w)_{\Omega_i} + (\nabla \cdot u_h, w)_{\Omega_i} = (f, w)_{\Omega_i} \quad \forall w \in W_{h,i}, \quad (5b)$$

$$\sum_{i=1}^n \langle u_{h,i}, n_i, \mu \rangle_{\Gamma_i} = 0 \quad \forall \mu \in M_H, \quad (5c)$$

In this formulation the pressure continuity (2e) is modeled via the mortar pressure function $\lambda_H$, while the flux continuity (2f) is imposed weakly on the coarse scale via (5c). For the above method to be well posed, the two scales must be chosen such that the mortar space is not too rich compared to the normal traces of the subdomain velocity spaces.

**Assumption 2.1.** Assume there exists a constant $C$ independent of $h$ and $H$ such that

$$\|\mu\|_{\Gamma_{i,j}} \leq C(\|Q_{h,i} \mu\|_{\Gamma_{i,j}} + \|Q_{h,j} \mu\|_{\Gamma_{i,j}}), \quad \forall \mu \in M_H, \quad 1 \leq i < j \leq n, \quad (6)$$

where $Q_{h,i} : L^2(\Gamma_i) \to V_{h,i} \cdot n_i|\Gamma_i$ is the $L^2$-projection operator from the mortar space onto the normal trace of the velocity space on subdomain $i$, i.e. for any $\phi \in L^2(\Gamma_i)$,

$$\langle \phi - Q_{h,i} \phi, v \cdot n_i \rangle_{\Gamma_i} = 0, \quad \forall v \in V_{h,i}. \quad (7)$$

This condition can be easily satisfied in practice by restricting the size of $H$ from below (see e.g. [28, 6, 22]). Under the above assumption, method (5) is solvable, stable, and accurate [7]. The following result has been shown in [7].
Theorem 2.1. If Assumption 2.1 holds, then method (5) has a unique solution and there exists a positive constant $C$, independent of $h$ and $H$, such that

$$
\| \nabla \cdot (u - u_h) \| \leq C \sum_{i=1}^{n} \| \nabla \cdot u \|_{r, \Omega_i} h^r, \quad 0 \leq r \leq l + 1, \quad (8)
$$

$$
\| u - u_h \| \leq C \sum_{i=1}^{n} (\| p \|_{s+1/2, \Omega_i} H^{s-1/2} + \| u \|_{r, \Omega_i} h^r + \| u \|_{r+1/2, \Omega_i} h^r H^{1/2}), \quad 1 \leq r \leq k + 1, \quad 0 \leq s \leq m + 1. \quad (9)
$$

Furthermore, if the problem on $\Omega$ is $H^2$-regular, then

$$
\| \hat{p} - p_h \| \leq C \sum_{i=1}^{n} (\| p \|_{s+1/2, \Omega_i} H^{s+1/2} + \| \nabla \cdot u \|_{t, \Omega_i} h^t H \quad (10)
$$

$$
+ \| u \|_{r, \Omega_i} h^r H + \| u \|_{r+1/2, \Omega_i} h^r H^{3/2}),
$$

$$
\| p - p_h \| \leq C \sum_{i=1}^{n} \| p \|_{t, \Omega_i} h^t + \| \hat{p} - p_h \|, \quad (11)
$$

where $1 \leq r \leq k + 1$, $0 \leq s \leq m + 1$, and $0 \leq t \leq l + 1$ and $\hat{p}$ is the $L^2$-projection of $p$ onto $W_h$.

2.4 Interface Formulation

Following [15], we formulate (5) as an interface problem for the mortar pressure. We decompose the solution to (5) into two parts: $u_h = u^*_h(\lambda_H) + \overline{u}_h$ and $p_h = p^*_h(\lambda_H) + \overline{p}_h$. The first component $(u^*_h, p^*_h) \in V_h^0 \times W_h$ solves subdomain problems with zero source and boundary conditions, and has $\lambda_H$ as a Dirichlet boundary condition along $\Gamma$, i.e. for $i = 1, \ldots, n$

$$
(K^{-1}u^*_h, v)_{\Omega_i} - (p^*_h, \nabla \cdot v)_{\Omega_i} = -(v \cdot n_i, \lambda_H)_{\partial \Omega_i \cap \Gamma} \quad \forall v \in V_{h,i}^0, \quad (12a)
$$

$$
(\alpha p^*_h, w)_{\Omega_i} + (\nabla \cdot u^*_h, w)_{\Omega_i} = 0 \quad \forall w \in W_{h,i}. \quad (12b)
$$

The second component $(\overline{u}_h, \overline{p}_h) \in V_h^{g_N} \times W_h$ solves subdomain problems with source $f$, boundary conditions $g_D$ and $g_N$ on $\partial \Omega$, and zero Dirichlet boundary conditions along $\Gamma$, i.e. for $i = 1, \ldots, n$

$$
(K^{-1}\overline{u}_h, v)_{\Omega_i} - (\overline{p}_h, \nabla \cdot v)_{\Omega_i} = -(v \cdot n_i, g_D)_{\partial \Omega_i \cap \Gamma_D} \quad \forall v \in V_{h,i}^0, \quad (13a)
$$

$$
(\alpha \overline{p}_h, w)_{\Omega_i} + (\nabla \cdot \overline{u}_h, w)_{\Omega_i} = (f, w)_{\Omega_i} \quad \forall w \in W_{h,i}. \quad (13b)
$$

Since the sum of (12a)-(12b) and (13a)-(13b) gives (5a)-(5b), all that remains to do is enforce equation (5c). Thus, the variational interface problem is to find $\lambda_H \in M_H$ such that

$$
\sum_{i=1}^{n} (-u^*_h(\lambda_H) \cdot n_i, \mu)_{\Gamma_i} = \sum_{i=1}^{n} (\overline{u}_h, n_i, \mu)_{\Gamma_i}, \quad \forall \mu \in M_H. \quad (14)
$$
Equivalently, we define bilinear forms $b_{H,i} : L^2(\Gamma_i) \times L^2(\Gamma_i) \to \mathbb{R}$ and $b_H : L^2(\Gamma) \times L^2(\Gamma) \to \mathbb{R}$ and a linear functional $g_H : L^2(\Gamma) \to \mathbb{R}$ by

$$b_{H,i}(\lambda_{H,i}, \mu) = \langle -u^*_{h,i}(\lambda_{H,i}) \cdot n_i, \mu \rangle_{\Gamma_i}, \quad (15a)$$

$$b_H(\lambda_H, \mu) = \sum_{i=1}^n b_{H,i}(\lambda_{H,i}, \mu), \quad (15b)$$

$$g_H(\mu) = \sum_{i=1}^n \langle \bar{u}_{h,i} \cdot n_i, \mu \rangle_{\Gamma_i}. \quad (15c)$$

With these definitions, (14) is equivalent to finding $\lambda_H \in M_H$ such that $b_H(\lambda_H, \mu) = g_H(\mu)$, for all $\mu \in M_H$. The distinction is made between bilinear forms (15a) and (15b) because the former measures the total flux across interface $\Gamma_i$ and requires no interprocessor communication, while the latter measures the total jump in flux across the set of all interfaces $\Gamma$ and hence does require interprocessor communication.

### 2.5 Interface Iteration

It is easy to check that $b_H$ is symmetric and positive semi-definite on $L^2(\Gamma)$. Moreover, it is positive definite on $M_H$ if Assumption 2.1 holds and either $\Gamma_D \neq \emptyset$ or $\alpha > 0$ [6, 7]. Therefore we solve the resulting discrete system with a Conjugate Gradient (CG) algorithm. Define linear operators $B_{H,i} : M_{H,i} \to M_{H,i}$ and $B_H : M_H \to M_H$ and a vector $g_H \in M_H$ corresponding to equations (15) by

$$\langle B_{H,i}\lambda_{H,i}, \mu \rangle_{\Gamma_i} = -\langle u^*_{h,i}(\lambda_{H,i}) \cdot n_i, \mu \rangle_{\Gamma_i} \forall \mu \in M_{H,i}, \quad (16a)$$

$$B_H \lambda_H = \sum_{i=1}^n B_{H,i} \lambda_{H,i}, \quad (16b)$$

$$\langle g_H, \mu \rangle_{\Gamma} = \sum_{i=1}^n \langle \bar{u}_{h,i} \cdot n_i, \mu \rangle_{\Gamma_i} \forall \mu \in M_H. \quad (16c)$$

Let $Q_{h,i}^T : V_{h,i} \cdot n_i|_{\Gamma_i} \to M_{H,i}$ be the $L^2$-orthogonal projection from the normal trace of the velocity space onto the mortar space. Note that (16a) and (16c) imply, respectively,

$$B_{H,i} = -Q_{h,i}^T u_{h,i}^*(\lambda_{H,i}) \cdot n_i, \quad g_H = \sum_{i=1}^n Q_{h,i}^T \bar{u}_{h,i} \cdot n_i. \quad (17)$$

Using this notation, the interface formulation is to find $\lambda_H \in M_H$ such that $B_H \lambda_H = g_H$. The operator $B_H$ is known as the Steklov-Poincaré operator [24].

Starting from an initial guess, we iterate on the value of $\lambda_H$ using the CG algorithm. On each CG iteration, we must evaluate the action of $B_H$ on $\lambda_H$. This is done with the following steps:
1. Project mortar data onto subdomain boundaries:

\[ \lambda_{H,i} \rightarrow \gamma_i. \]

2. Solve the set of subdomain problems (12) with Dirichlet boundary data \( \gamma_i \).

3. Project the resulting fluxes from Step 2 back onto the mortar space:

\[ -u^*_h(\gamma_i) \cdot n_i \rightarrow -Q^T_{h,i} u^*_h(\gamma_i) \cdot n_i. \]

4. Compute flux jumps across all subdomain interfaces \( \Gamma_{i,j} \):

\[ B_{H,i} \lambda_{H,i} = - \sum_{i=1}^{n} Q^T_{h,i} u^*_h(\gamma_i) \cdot n_i. \]

Steps 1-3 evaluate the action of the flux operator \( B_{H,i} \) as in (17) and are done by every subdomain in parallel. Step 4 evaluates the action of the jump operator \( B_{H} \) as in (16b) and requires interprocessor communication across every subdomain interface.

### 3 Multiscale Flux Basis Implementation

Observe that the dominant computational cost in each CG iteration is in the evaluation of the flux operator \( B_{H,i} \), which requires solving one Dirichlet subdomain problem per subdomain. One way to potentially reduce this computational cost is with the following approach.

Before the CG algorithm begins, compute and store the flux responses associated with each mortar degree of freedom, on each subdomain independently.

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This is what we call the Multiscale Flux Basis. Its assembly requires solving only a fixed number of Dirichlet subdomain problems (12), see Figure 1. After these
solves are completed, the action of $B_{H,i}$ is reduced to taking a linear combination of Multiscale Flux Basis functions. Therefore, if the number of CG iterations exceeds the maximum number of mortar degrees of freedom on any subdomain, then the computational cost will be reduced in terms of fewer required subdomain solves, and should yield faster runtime.

### 3.1 Assembly of the Multiscale Flux Basis

To compute each function in the Multiscale Flux Basis, we shall apply Steps 1-3 from the interface iteration in order to evaluate the action of the operator $B_{H,i}$ on each mortar basis function, on each subdomain independently. Let there be exactly $N_{H,i}$ mortar degrees of freedom on subdomain $i$ and define $\{\phi_{H,i}^{(k)}\}_{k=1}^{N_{H,i}}$ to be the mortar basis functions for $M_{H,i}$. Then for $\lambda_{H,i} \in M_{H,i}$ we may express

$$\lambda_{H,i} = \sum_{k=1}^{N_{H,i}} \lambda_{H,i}^{(k)} \phi_{H,i}^{(k)}.$$

Consider the $k$-th mortar basis function $\phi_{H,i}^{(k)}$. Computing the Multiscale Flux Basis function corresponding to $\phi_{H,i}^{(k)}$ involves the following steps.

1. Project this mortar basis function onto the subdomain boundary:

$$Q_{h,i} \phi_{H,i}^{(k)} = \gamma_{i}^{(k)}.$$

2. Solve problem (12) on each subdomain $\Omega_i$ with Dirichlet interface condition $\gamma_{i}^{(k)}$, i.e. find $u_h^* = u_h^*(\gamma_{i}^{(k)})$ and $p_h^* = p_h^*(\gamma_{i}^{(k)})$ such that

$$(K^{-1} u_h^*, v)_{\Omega_i} - (p_h^*, \nabla \cdot v)_{\Omega_i} = - (v \cdot n_i, \gamma_{i}^{(k)})_{\partial \Omega_i \cap \Gamma} \quad \forall v \in V_{h,i}^0,$$

$$(\alpha p_h^*, w)_{\Omega_i} + (\nabla \cdot u_h^*, w)_{\Omega_i} = 0 \quad \forall w \in W_{h,i}.$$

3. Project the resulting boundary flux back onto the mortar space:

$$\psi_{H,i}^{(k)} = - Q_{h,i}^T u_h^*(\gamma_{i}^{(k)}) \cdot n_i.$$

Repeating this procedure for $k = 1, \ldots, N_{H,i}$ forms the Multiscale Flux Basis for subdomain $\Omega_i$:

$$\{\psi_{H,i}^{(1)}, \psi_{H,i}^{(2)}, \ldots, \psi_{H,i}^{(N_{H,i})}\} \subset M_{H,i}.$$

**Remark 3.1.** Note that each mortar basis function $\phi_{H,i}^{(k)}$ on interface $\Gamma_{i,j}$ corresponds to exactly two different Multiscale Flux Basis functions, one for $\Omega_i$ and one for $\Omega_j$. 

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Remark 3.2. Whereas the original MMMFEM implementation requires each processor to perform the exact same number of subdomain solves during the interface iteration process, our implementation may have each processor perform a different number of subdomain solves in assembling its Multiscale Flux Basis. This is because there may be a varying number of degrees of freedom in each mortar $M_{H,i,j}$ and subdomains may share portions of their boundaries with $\Gamma_D$ and $\Gamma_N$.

Remark 3.3. The extra storage cost in saving the Multiscale Flux Basis functions is equal to the square of the size of the mortar space on each subdomain, $N_H^2$. This is because each multiscale flux basis function belongs to the mortar space $M_{H,i}$ of dimension $N_{H,i}$, and there are exactly $N_{H,i}$ basis functions to be computed. Since only interface data needs to be stored, the storage cost is significantly lower than the storage cost for the variational multiscale method and multiscale finite elements, where the basis functions are defined on the entire local fine grid.

3.2 Using the Multiscale Flux Basis in the Interface Iteration

To use the Multiscale Flux Basis to replace Steps 1-3 in the interface iteration, we need only observe that the flux operator $B_{H,i}$ is linear. Therefore,

$$B_{H,i}(\lambda_{H,i}) = B_{H,i} \left( \sum_{k=1}^{N_{H,i}} \lambda_{H,i}(k) \phi_{H,i}(k) \right) = \sum_{k=1}^{N_{H,i}} \lambda_{H,i}(k) B_{H,i} \left( \phi_{H,i}(k) \right) = \sum_{k=1}^{N_{H,i}} \lambda_{H,i}(k) \psi_{H,i}(k).$$ (19)

In other words, to compute the resulting flux on subdomain $\Omega_i$ from Dirichlet data $\lambda_{H,i}$, we simply take a linear combination of the Multiscale Flux Basis functions $\psi_{H,i}(k)$ using the same scalars which express $\lambda_{H,i}$ in terms of its mortar basis functions $\phi_{H,i}(k)$. This demonstrates the equivalence of the original MMMFEM implementation to our new Multiscale Flux Basis implementation.

Remark 3.4. In the original MMMFEM implementation, fine scale pressure and velocity variables may also be updated iteratively in the interface iteration. In the new Multiscale Flux Basis implementation, this convention should be dropped, because storing arrays of these fine scale variables for each mortar degree of freedom would be an unnecessary burden on memory. Instead, we perform one additional Dirichlet subdomain solve after the CG iteration has converged in order to recover the fine scale pressure and velocity.

4 Numerical Examples

The algorithm described in the previous section was implemented in the parallel flow simulator PARCEL, which is programmed in FORTRAN. The domain decomposition uses spatially conforming rectangular subdomains (in 2-D) or brick subdomains...
Within each of these subdomains, the fine grid is comprised of the lowest order \textit{Raviart-Thomas mixed finite element spaces} on rectangles (in 2-D) or bricks (in 3-D) \cite{25, 21}, which are allowed to be spatially non-conforming across subdomain interfaces. On these subdomain interfaces, a coarse grid is comprised of continuous or discontinuous, linear or quadratic mortar spaces. Three numerical methods are compared:

\textbf{Method 1.} \textit{Original MMMFEM implementation, no interface preconditioner.}

\textbf{Method 2.} \textit{Original MMMFEM implementation, Balancing preconditioner.}

\textbf{Method 3.} \textit{New Multiscale Flux Basis implementation, no preconditioner.}

Unless otherwise noted, the tolerance for the relative residual in the \textit{CG} algorithm is taken to be 1e-06.

\textbf{Remark 4.1.} The balancing preconditioner used in the tests has been described in \cite{13, 22, 7}. It involves solving Neumann subdomain problems and a coarse problem which provides global exchange of information across subdomains. This causes the condition number of the interface problem to grow more modestly versus non-preconditioned \textit{CG} as the grids are refined or the number of subdomains increases. The cost for one preconditioned iteration is three subdomain solves and two coarse solves.

Four example problems are considered: a 2-D problem with smooth permeability, a 2-D problem with a rough permeability, a 3-D problem with smooth permeability, and a 2-D problem with adaptive mesh refinement. In the first three examples we solve each problem using a fixed fine grid several times. Each time we increase the number of subdomains, \textit{i.e.}, refine the coarse grid. This causes the interface problem to become larger and more ill-conditioned, hence increasing the number of \textit{CG} iterations. Tables are provided which compare both the number of \textit{CG} iterations and maximum number of subdomain solves required by the three methods\footnote{Recall Remark 3.2.}. In this way, the new multiscale flux basis implementation can be directly compared to the original MMMFEM implementation. No error norms are reported in these tests, because all three methods produce the same solution within roundoff error.

For the first two examples we also provide tests comparing the accuracy and the cost of the MMMFEM solution to a fine scale solution.

The fourth example involves adaptive mesh refinement and illustrates the greater flexibility of the MMMFEM compared to existing multiscale methods. It also shows that the gain in efficiency from the new implementation is increased when grid adaptivity is employed.

\textbf{Remark 4.2.} It should be noted that under a fixed fine grid, as the number of subdomains is increased, the size of the local subdomain problems becomes smaller.
4.1 Example 1: 2-D problem with a smooth solution

This example is a 2-D problem on the domain $\Omega = (0, 1)^2$ with a fixed global fine grid of $120 \times 120$ elements. The solution is given by $p(x, y) = x^3y^4 + x^2 + \sin(xy)\cos(y)$, and the coefficient $K$ is a smooth, full tensor defined by

$$K = \begin{pmatrix}
(x + 1)^2 + y^2 & \sin(xy) \\
\sin(xy) & (x + 1)^2
\end{pmatrix}.$$ 

Boundaries $\{y = 0\}$ and $\{y = 1\}$ are Dirichlet type and boundaries $\{x = 0\}$ and $\{x = 1\}$ are Neumann type.

Table 1: Example 1 using continuous linear mortars with 3 elements per edge.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>Method 1</th>
<th>Method 2</th>
<th>Method 3</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
</tr>
<tr>
<td>2 $\times$ 2 = 4</td>
<td>14</td>
<td>17*</td>
<td>11</td>
</tr>
<tr>
<td>3 $\times$ 3 = 9</td>
<td>29</td>
<td>32*</td>
<td>19</td>
</tr>
<tr>
<td>4 $\times$ 4 = 16</td>
<td>42</td>
<td>45</td>
<td>24</td>
</tr>
<tr>
<td>5 $\times$ 5 = 25</td>
<td>54</td>
<td>57</td>
<td>26</td>
</tr>
<tr>
<td>6 $\times$ 6 = 36</td>
<td>65</td>
<td>68</td>
<td>27</td>
</tr>
<tr>
<td>7 $\times$ 7 = 49</td>
<td>75</td>
<td>78</td>
<td>26</td>
</tr>
<tr>
<td>8 $\times$ 8 = 64</td>
<td>86</td>
<td>89</td>
<td>26</td>
</tr>
</tbody>
</table>

* - denotes fewest number of solves

Table 1 shows results for Example 1 using continuous linear mortars with 3 elements per edge. Method 1 performs the best until we reach the $4 \times 4$ subdomain case. After this point Method 3 becomes the most efficient in terms of subdomain solves. Observe that the number of CG iterations increases with the number of subdomains, and that Method 1 requires one subdomain solve per iteration plus three additional subdomain solves. The balancing preconditioner used in Method 2 causes the number of CG iterations to grow more modestly with the number of subdomains, but this method is still more costly in terms of subdomain solves. Notice that Method 3 always requires a fixed maximum number of subdomain solves (except in the $2 \times 2$ case where only two out of four edges of each subdomain have mortars). This table demonstrates that as the number of subdomains is increased, there is a point after which Method 3 performs best. We found this to be the case for most tests we ran.

Remark 4.3. Recall that the Balancing preconditioner involves two additional coarse grid solves per CG iteration. Thus even in cases where Method 2 required fewer subdomain solves, Method 3 was more efficient in terms of CPU time, as the time for the coarse solves was not negligible. We do not report CPU times in this paper, since they depend on the particular implementation of the coarse solve in the Balancing preconditioner.
Table 2: Example 1 using continuous quadratic mortars with 2 elements per interface.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>Method 1</th>
<th></th>
<th>Method 2</th>
<th></th>
<th>Method 3</th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
<td>Solves</td>
</tr>
<tr>
<td>2 × 2 = 4</td>
<td>16</td>
<td>19*</td>
<td>12</td>
<td>44</td>
<td>16</td>
<td>53</td>
</tr>
<tr>
<td>3 × 3 = 9</td>
<td>35</td>
<td>38*</td>
<td>17</td>
<td>61</td>
<td>34</td>
<td>63</td>
</tr>
<tr>
<td>4 × 4 = 16</td>
<td>51</td>
<td>54*</td>
<td>20</td>
<td>70</td>
<td>51</td>
<td>63</td>
</tr>
<tr>
<td>5 × 5 = 25</td>
<td>65</td>
<td>68</td>
<td>21</td>
<td>73</td>
<td>65</td>
<td>63*</td>
</tr>
<tr>
<td>6 × 6 = 36</td>
<td>78</td>
<td>81</td>
<td>22</td>
<td>76</td>
<td>78</td>
<td>63*</td>
</tr>
<tr>
<td>7 × 7 = 49</td>
<td>91</td>
<td>94</td>
<td>21</td>
<td>73</td>
<td>91</td>
<td>63*</td>
</tr>
<tr>
<td>8 × 8 = 64</td>
<td>103</td>
<td>107</td>
<td>21</td>
<td>73</td>
<td>103</td>
<td>63*</td>
</tr>
</tbody>
</table>

* - denotes fewest number of solves

In Table 2 we report results for Example 1 with continuous quadratic mortars with 2 elements per edge. This slightly increases the required work for Method 1 and slightly decreases the work for Method 2. However, for Method 3 this change nearly doubles the amount of subdomain solves required due to the increase in mortar degrees of freedom per subdomain. This means that initially our method solves more subdomain problems than the other two, and the computational efficiency of Method 3 is not observed until the 5 × 5 case. This difference versus the previous table shows that the number of mortar degrees of freedom per subdomain is an important parameter which determines the relative computationally efficiency of the Multiscale Flux Basis implementation.

Table 3: Numerical error and computational cost for the fine scale solution in Example 1.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>pres-L2-err</th>
<th>vel-L2-err</th>
<th>CGIter</th>
<th>Solves</th>
</tr>
</thead>
<tbody>
<tr>
<td>2 × 2 = 4</td>
<td>7.1657E-05</td>
<td>7.1848E-05</td>
<td>58</td>
<td>123</td>
</tr>
<tr>
<td>3 × 3 = 9</td>
<td>7.1955E-05</td>
<td>7.9269E-05</td>
<td>72</td>
<td>163</td>
</tr>
<tr>
<td>4 × 4 = 16</td>
<td>7.1968E-05</td>
<td>8.7311E-05</td>
<td>85</td>
<td>123</td>
</tr>
<tr>
<td>5 × 5 = 25</td>
<td>7.2211E-05</td>
<td>9.5265E-05</td>
<td>96</td>
<td>99</td>
</tr>
<tr>
<td>6 × 6 = 36</td>
<td>7.2260E-05</td>
<td>1.0281E-04</td>
<td>107</td>
<td>83</td>
</tr>
</tbody>
</table>

To illustrate the accuracy of the MMMFEM and the efficiency of the proposed new implementation, we compare the quality and cost of the multiscale solution to these of the fine scale solution. The latter is computed using the same domain decomposition algorithm with Method 3, but with fine scale Lagrange multipliers. In Table 3 we report the pressure and velocity errors and the cost of the interface iteration. This type of test is comparable to a standard mixed finite element algorithm without domain decomposition. Indeed, the recorded error norms remain nearly constant as the number of subdomains is increased.
In comparison, Table 4 shows results for the MMMFEM using Method 3 with linear mortars and a single element per interface. This subdomain configuration is very much akin to the variational multiscale methods and multiscale finite element methods mentioned in the introduction. We note that the MMMFEM requires significantly smaller number of subdomain solves, while at the same time resolves the flow very well, as seen in Figure 2 where a comparison of the plots of the computed fine scale and multiscale solutions with $5 \times 5$ subdomains is shown. The error norms reported in Table 4 indicate that the error is larger for the multiscale solution, but does decrease as the number of subdomain is increased. Figure 3 shows that the locations with greatest error in the multiscale solution are along the subdomain interfaces.

Table 4: Numerical error and computational cost for the multiscale solution using Method 3 with a single linear mortar per interface in Example 1.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>pres-L2-err</th>
<th>vel-L2-err</th>
<th>CGIter</th>
<th>Solves</th>
</tr>
</thead>
<tbody>
<tr>
<td>$2 \times 2 = 4$</td>
<td>1.2966E-02</td>
<td>4.4386E-02</td>
<td>8</td>
<td>11</td>
</tr>
<tr>
<td>$3 \times 3 = 9$</td>
<td>7.1036E-03</td>
<td>3.6534E-02</td>
<td>22</td>
<td>19</td>
</tr>
<tr>
<td>$4 \times 4 = 16$</td>
<td>4.2496E-03</td>
<td>3.0038E-02</td>
<td>33</td>
<td>19</td>
</tr>
<tr>
<td>$5 \times 5 = 25$</td>
<td>2.7673E-03</td>
<td>2.5191E-02</td>
<td>42</td>
<td>19</td>
</tr>
<tr>
<td>$6 \times 6 = 36$</td>
<td>1.9159E-03</td>
<td>2.1527E-02</td>
<td>51</td>
<td>19</td>
</tr>
</tbody>
</table>

Figure 2: Computed pressure (color) and velocity (arrows) in Example 1: fine scale solution (left) and multiscale solution with a single linear mortar per interface (right).
Figure 3: Pressure error (left) and magnitude of velocity error (right) for the multiscale solution with a single linear mortar per interface in Example 1.

4.2 Example 2: 2-D problem with rough heterogeneous permeability

This problem uses a 2-D heterogeneous permeability field, obtained from the Society of Petroleum Engineers (SPE) Comparative Solution Project\(^2\). The domain is \(\Omega = (0, 60) \times (0, 220)\) with a fixed global fine grid of 60 \(\times\) 220 elements. Pressure values of one and zero are specified on the left and right boundaries, respectively. No flow is specified on the top and bottom boundaries. A plot of the permeability field is shown on the left in Figure 4.

Table 5: Example 2 using continuous linear mortars with 2 elements per interface.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>Method 1</th>
<th>Method 2</th>
<th>Method 3</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
</tr>
<tr>
<td>2 (\times) 2 = 4</td>
<td>13</td>
<td>15</td>
<td>8</td>
</tr>
<tr>
<td>3 (\times) 2 = 6</td>
<td>19</td>
<td>21</td>
<td>15</td>
</tr>
<tr>
<td>2 (\times) 4 = 8</td>
<td>25</td>
<td>27</td>
<td>18</td>
</tr>
<tr>
<td>2 (\times) 5 = 10</td>
<td>37</td>
<td>39</td>
<td>29</td>
</tr>
<tr>
<td>3 (\times) 4 = 12</td>
<td>37</td>
<td>39</td>
<td>28</td>
</tr>
<tr>
<td>3 (\times) 5 = 15</td>
<td>51</td>
<td>53</td>
<td>37</td>
</tr>
</tbody>
</table>

* - denotes fewest number of solves

Table 5 shows the results for Example 2 using continuous linear mortars with 2 elements per edge. Method 3 requires at most 26 solves per subdomain and is compu-

\(^2\)For more information, see http://www.spe.org/csp
Figure 4: Example 2: permeability field (left), fine scale solution (middle), and multiscale solution with $3 \times 5$ subdomains and a single linear mortar per interface (right).

The computational efficiency of Method 3 is demonstrated in Figure 4. It is observed that Method 3 is computationally more efficient than Methods 1 and 2 for all subdomain configurations. As the number of subdomains is increased, the improvement of Method 3 over Methods 1 and 2 becomes greater.

A comparison between the fine scale solution and the multiscale solution with $3 \times 5$ subdomains is presented in Figure 4. We observe a very good match between the two solutions. We note that the number of subdomain solves required by Method 3 for the multiscale solution, 26, is significantly less than Methods 1-3 used for computing the fine scale solution, which require 388, 84, and 130 subdomain solves, respectively.

Table 6 shows the results for Example 2 using continuous quadratic mortars with 2 elements per interface. Compared to the previous table, the increased number of mortar degrees of freedom per interface leads to more subdomain solves for Method 3, the maximum number being 62. Nevertheless, Method 3 is still more computationally efficient than Methods 1 and 2 for 10 and more subdomains.

### 4.3 Example 3: 3-D problem with a smooth solution

This example is a 3-D problem on the domain $\Omega = (0, 1)^3$ with a fixed global fine grid of $48 \times 48 \times 48$ elements. The solution is given by $p(x, y, z) = x + y + z - 1.5$, and
Table 6: Example 2 using continuous quadratic mortars with 2 elements per interface.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>Method 1</th>
<th>Method 2</th>
<th>Method 3</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
</tr>
<tr>
<td>2 × 2 = 4</td>
<td>17</td>
<td>19*</td>
<td>15</td>
</tr>
<tr>
<td>3 × 2 = 6</td>
<td>30</td>
<td>32*</td>
<td>23</td>
</tr>
<tr>
<td>2 × 4 = 8</td>
<td>39</td>
<td>41*</td>
<td>25</td>
</tr>
<tr>
<td>2 × 5 = 10</td>
<td>56</td>
<td>58</td>
<td>39</td>
</tr>
<tr>
<td>3 × 4 = 12</td>
<td>53</td>
<td>55</td>
<td>33</td>
</tr>
<tr>
<td>3 × 5 = 15</td>
<td>92</td>
<td>94</td>
<td>46</td>
</tr>
</tbody>
</table>

* - denotes fewest number of solves

the coefficient \( K \) is a smooth full tensor defined by

\[
K = \begin{pmatrix}
x^2 + y^2 + 1 & 0 & 0 \\
0 & z^2 + 1 & \sin(xy) \\
0 & \sin(xy) & x^2y^2 + 1
\end{pmatrix}.
\]

Boundaries \( \{ y = 0 \} \) and \( \{ y = 1 \} \) are Dirichlet type and the rest of the boundary is Neumann type.

Figure 5 shows the computed multiscale solution and its error for Example 3 with \( 4 \times 4 \times 4 \) subdomains and a single linear mortar per interface. Table 7 shows the computational cost for Methods 1-3 with various coarse grids. Method 3 requires at most 27 solves per subdomain and outperforms Methods 1 and 2 for all subdomain configurations.

Table 7: Example 3 using linear mortars with one element per interface.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>Method 1</th>
<th>Method 2</th>
<th>Method 3</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
</tr>
<tr>
<td>2 × 2 × 2  = 8</td>
<td>28</td>
<td>31</td>
<td>11</td>
</tr>
<tr>
<td>2 × 2 × 3  = 12</td>
<td>33</td>
<td>36</td>
<td>12</td>
</tr>
<tr>
<td>2 × 3 × 3  = 18</td>
<td>37</td>
<td>40</td>
<td>13</td>
</tr>
<tr>
<td>3 × 3 × 3  = 27</td>
<td>46</td>
<td>49</td>
<td>13</td>
</tr>
<tr>
<td>3 × 3 × 4  = 36</td>
<td>50</td>
<td>53</td>
<td>13</td>
</tr>
<tr>
<td>3 × 4 × 4  = 48</td>
<td>55</td>
<td>58</td>
<td>13</td>
</tr>
<tr>
<td>4 × 4 × 4  = 64</td>
<td>60</td>
<td>63</td>
<td>13</td>
</tr>
</tbody>
</table>

* - denotes fewest number of solves

Table 8 shows the results for Example 3 using quadratic mortars with one element per interface with the usual relative residual CG tolerance of 1e-06. Method 3 requires at most 57 solves per subdomain. It is the fastest method on coarser domain decompositions, but Method 2 outperforms it slightly on 27 or more subdomains.
Figure 5: Example 3: computed multiscale solution (left) and its error (right) on \(4 \times 4 \times 4\) subdomains with a single linear mortar per interface.

Table 8: Example 3 using quadratic mortars with one element per interface. Relative residual CG tolerance = 1e-06.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>Method 1</th>
<th>Method 2</th>
<th>Method 3</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
</tr>
<tr>
<td>(2 \times 2 \times 2 = 8)</td>
<td>36</td>
<td>39</td>
<td>13</td>
</tr>
<tr>
<td>(2 \times 2 \times 3 = 12)</td>
<td>41</td>
<td>44</td>
<td>13</td>
</tr>
<tr>
<td>(2 \times 3 \times 3 = 18)</td>
<td>47</td>
<td>50</td>
<td>14</td>
</tr>
<tr>
<td>(3 \times 3 \times 3 = 27)</td>
<td>56</td>
<td>59</td>
<td>14</td>
</tr>
<tr>
<td>(3 \times 3 \times 4 = 36)</td>
<td>60</td>
<td>63</td>
<td>14</td>
</tr>
<tr>
<td>(3 \times 4 \times 4 = 48)</td>
<td>64</td>
<td>67</td>
<td>14</td>
</tr>
<tr>
<td>(4 \times 4 \times 4 = 64)</td>
<td>69</td>
<td>72</td>
<td>14</td>
</tr>
</tbody>
</table>

* - denotes fewest number of solves

When a tighter tolerance is imposed on the CG on the interface, all three methods perform more CG iterations. Under Methods 1 and 2, this also requires performing more subdomain solves. For Method 3, however, the maximum number of solves per subdomain is unaffected by this change in tolerance. This is illustrated in Table 9, which shows the results for relative residual CG tolerance of 1e-09. In this case Method 3 is the most computationally efficient for all subdomain configurations.
Table 9: Example 3 using quadratic mortars with one element per interface. Relative residual CG tolerance = 1e-09.

<table>
<thead>
<tr>
<th>Subdomains</th>
<th>Method 1</th>
<th></th>
<th>Method 2</th>
<th></th>
<th>Method 3</th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
<td>Solves</td>
<td>CGIter</td>
<td>Solves</td>
</tr>
<tr>
<td>$2 \times 2 = 8$</td>
<td>48</td>
<td>51</td>
<td>19</td>
<td>66</td>
<td>48</td>
<td>30*</td>
</tr>
<tr>
<td>$2 \times 2 \times 2 = 8$</td>
<td>48</td>
<td>51</td>
<td>19</td>
<td>66</td>
<td>48</td>
<td>30*</td>
</tr>
<tr>
<td>$2 \times 3 = 12$</td>
<td>56</td>
<td>59</td>
<td>19</td>
<td>67</td>
<td>56</td>
<td>39*</td>
</tr>
<tr>
<td>$2 \times 3 \times 3 = 12$</td>
<td>56</td>
<td>59</td>
<td>19</td>
<td>67</td>
<td>56</td>
<td>39*</td>
</tr>
<tr>
<td>$3 \times 3 = 27$</td>
<td>74</td>
<td>77</td>
<td>20</td>
<td>72</td>
<td>74</td>
<td>57*</td>
</tr>
<tr>
<td>$3 \times 3 \times 3 = 27$</td>
<td>74</td>
<td>77</td>
<td>20</td>
<td>72</td>
<td>74</td>
<td>57*</td>
</tr>
<tr>
<td>$3 \times 4 = 36$</td>
<td>79</td>
<td>82</td>
<td>21</td>
<td>75</td>
<td>79</td>
<td>57*</td>
</tr>
<tr>
<td>$3 \times 4 \times 3 = 36$</td>
<td>79</td>
<td>82</td>
<td>21</td>
<td>75</td>
<td>79</td>
<td>57*</td>
</tr>
<tr>
<td>$3 \times 4 \times 4 = 48$</td>
<td>84</td>
<td>87</td>
<td>21</td>
<td>75</td>
<td>85</td>
<td>57*</td>
</tr>
<tr>
<td>$4 \times 4 = 64$</td>
<td>92</td>
<td>95</td>
<td>21</td>
<td>75</td>
<td>92</td>
<td>57*</td>
</tr>
</tbody>
</table>

* - denotes fewest number of solves

4.4 Example 4: Mesh Adaptivity

The final example shows another benefit to the MMMFEM, the utilization of adaptive mesh refinement. The calculation of a posteriori error indicators within each subdomain allows the MMMFEM to refine only those subdomains where the error is highest. Mortars which touch refined subdomains are also refined to maintain accuracy. For more details, see [27, 7].

The permeability $K$ is a single realization of a stochastic permeability field on the domain $(0,1)^2$. A Karhunen-Loève expansion for the log permeability $Y = \ln(K)$ (a scalar quantity) is computed from the specified covariance function

$$C_Y(x, \bar{x}) = \sigma_Y^2 \exp \left[ -\frac{|x_1 - \bar{x}_1|}{\eta_1} - \frac{|x_2 - \bar{x}_2|}{\eta_2} \right].$$

The parameters used for this test are correlation lengths $\eta_1 = 0.25$, $\eta_2 = 0.125$, and variance $\sigma_Y = 2.1$. The series was truncated after 400 terms. For the exact procedure on computing the eigenfunctions and eigenvalues of this series semi-analytically, the interested reader can consult Appendix A in [29].

This test was performed on $5 \times 5 = 25$ subdomains, initially starting with $2 \times 2$ subdomain grids and continuous linear mortars with 1 element per edge. The permeability field and its corresponding solution on the fourth level of mesh refinement are shown in Figure 6. Using Method 1, each subdomain performed 283 subdomain solves, roughly one for each CG iteration on each of the 4 grid levels. Using Method 3, the number of subdomain solves after 4 levels of mesh refinement is shown in the figure on top of the permeability plot. The maximum number of subdomain solves is 160 and the minimum number is 56.

We can draw two conclusions from this example. First, since the computational savings of the multiscale flux basis implementation happen on each level of adaptive mesh refinement, the overall savings after all levels are complete is amplified by the
Second, the workload for each processor may become increasingly unbalanced due to a large variation in the number of mortar degrees of freedom per subdomain. Nevertheless, even if the algorithm is only as fast as its slowest processor, the Multiscale Flux Basis implementation is still faster than the original implementation. One can take full advantage of the computational efficiency of the new method in adaptive mesh refinement setting by implementing load balancing.

Figure 6: Permeability field for Example 4 on mesh refinement level 4 (left) and its corresponding solution (right). Numbers indicate the total number of subdomain solves using the Multiscale Flux Basis implementation.

5 Conclusions

In this paper we present a new implementation of the MMMFEM, which makes it comparable in computational cost to existing multiscale methods. The MMMFEM provides extra flexibility in the ability to vary locally and adaptively both the coarse scale and the fine scale grids. Moreover, the fine scale grids can be completely non-matching across coarse interfaces. The proposed implementation is based on precomputation of a Multiscale Flux Basis and using it to compute solutions to subdomain solves during the global coarse scale iteration.

The numerical examples demonstrate that the new implementation is more computationally efficient than the original implementation in many cases. The number of subdomain solves required for the construction of the Multiscale Flux Basis depends only on the number of mortar degrees of freedom per subdomain and not on the size of the global problem. Therefore the new implementation outperforms the original
one for large problems. Moreover, if the Multiscale Flux Basis implementation is used repeatedly, as in the case of adaptive mesh refinement, then the computational savings are amplified. Even greater computational gain is observed when this approach is combined with stochastic collocation for uncertainty quantification, which requires a large number of deterministic simulations. This extension will be discussed in a forthcoming paper.

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References


